

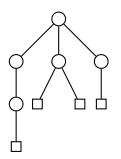




Combinatorial Model

•000

Remove all leaves!



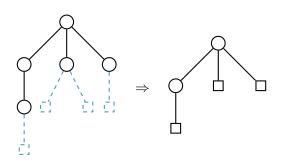




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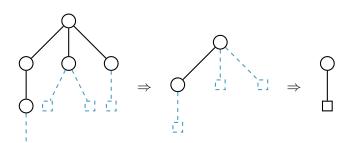




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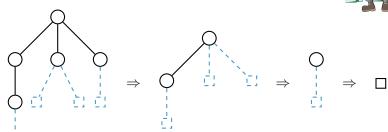


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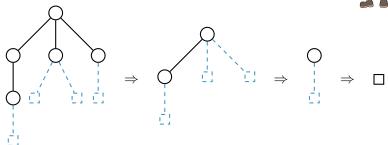


Combinatorial Model

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► Remove all leaves!





Parameters of Interest:

- ► Size of *r*th reduction
- ► Age: # of possible reductions



Reduction \rightarrow Expansion

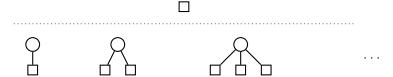
- modelling reduction directly: not suitable
- ▶ instead: see inverse operation, growing trees



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Combinatorial Model

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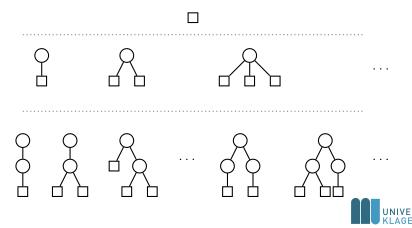




Reduction \rightarrow Expansion

Combinatorial Model

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- ▶ instead: see inverse operation, growing trees



Expansion operators

- \blacktriangleright F... family of plane trees; bivariate generating function f
- ightharpoonup expansion operator $\Phi \Rightarrow \Phi(f)$ counts expanded trees



Expansion operators

Combinatorial Model

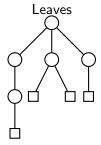
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Leaf expansion Φ

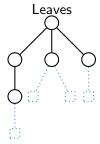
- inverse operation to leaf reduction
 - attach leaves to all current leaves (required)
 - attach leaves to inner nodes (optional)

$$\square \triangleq t, \ \bigcirc \triangleq z \quad \Rightarrow \quad \Phi(t) = zt + zt^2 + zt^3 + \cdots$$

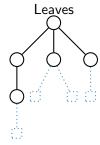




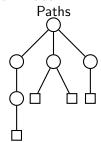




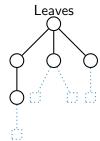




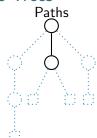
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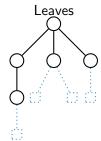




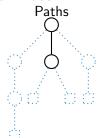
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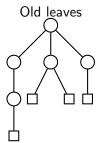




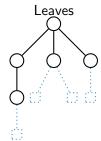


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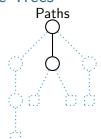


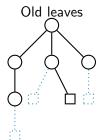




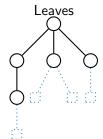


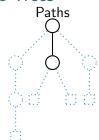
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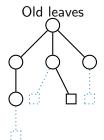


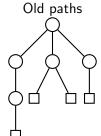




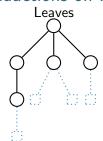




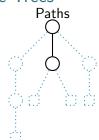


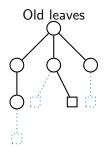


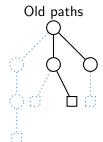




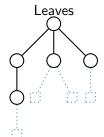
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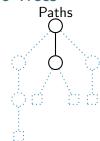


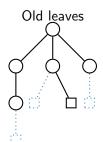


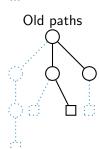


Combinatorial Model

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Parameters of Interest:

- tree size after r reductions
- cumulative reduction size



Proposition

► T...rooted plane trees



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- ▶ T(z,t)... BGF for $T(z \rightsquigarrow inner nodes, t \rightsquigarrow leaves)$



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$$\Rightarrow T(z,t) = \frac{1-(z-t)-\sqrt{1-2(z+t)+(z-t)^2}}{2}$$

Proof. Symbolic equation

$$T = \Box + T T \cdots T$$

translates into

$$T(z,t) = t + z \cdot \frac{T(z,t)}{1 - T(z,t)}$$



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$$\Phi(f(z,t)) = (1-t)f\left(\frac{z}{(1-t)^2}, \frac{zt}{(1-t)^2}\right)$$



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- **Expansion:**



► In total:

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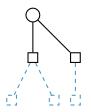
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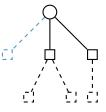


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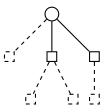




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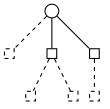
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University

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 \triangleright As Φ is linear, this proves the proposition.



▶ Functional equation: $T(z,t) = \Phi(T(z,t)) + t$



- Functional equation: $T(z,t) = \Phi(T(z,t)) + t$
- With $z = u/(1+u)^2$ and by some manipulations

$$\Phi^{r}(T(z,t))|_{t=z} = \frac{1 - u^{r+2}}{(1 - u^{r+1})(1 + u)} T\left(\frac{u(1 - u^{r+1})^{2}}{(1 - u^{r+2})^{2}}, \frac{u^{r+1}(1 - u)^{2}}{(1 - u^{r+2})^{2}}\right)$$



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- ▶ BGF $G_r(z, v)$ for size comparison: z tracks original size, v size of r-fold reduced tree
- Intuition: v "remembers" size while tree family is expanded

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SageMath Demo

https://benjamin-hackl.at/downloads/talks/2019-12-10-seminar-strobl.slides.html



Theorem (H.-Heuberger-Kropf-Prodinger)

After r reductions of a random tree of size n, the remaining size $X_{n,r}$ has mean and variance

$$\mathbb{E}X_{n,r} = \frac{n}{r+1} - \frac{r(r-1)}{6(r+1)} + O(n^{-1}),$$

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and $X_{n,r}$ is asymptotically normally distributed.



Cutting leaves

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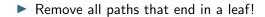
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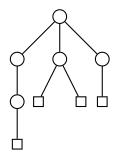
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Proof insights:

- $ightharpoonup \mathbb{E} X_{n,r}$ and $\mathbb{V} X_{n,r}$ follow via singularity analysis
- Asymptotic normality: $n X_{n,r}$ is a tree parameter with small toll function, limit law by Wagner (2015)

Pruning



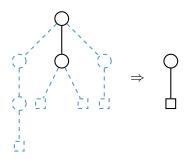






Pruning

▶ Remove all paths that end in a leaf!







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Proposition / Conj.: More than just equality in distribution

 \mathcal{T}_n ... trees of size n, ρ_P ... path reduction, ρ_L ... leaf reduction

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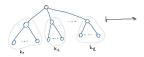
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Bijection for removed structures for r = 1:

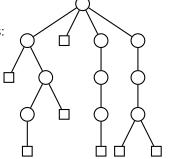




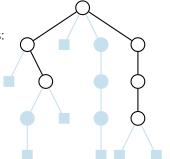


Work in progress for r > 2.

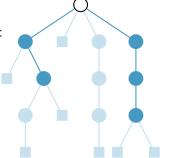




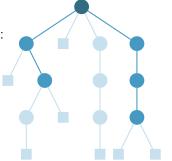






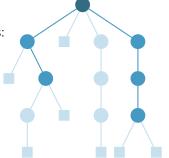






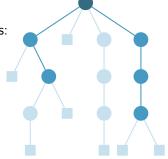


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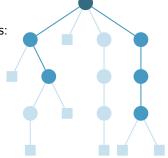


Observation

Total # of branches \triangleq # of leaves in all reduction stages



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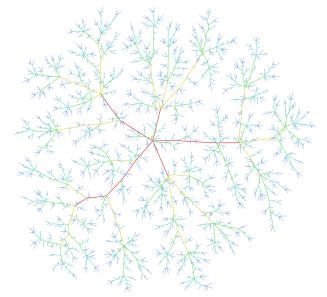
Observation

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Proof: all branches end in exactly one leaf (at some point).



Visualized Branches in a Tree





Branches in a Tree - Result

Theorem (H.-Heuberger-Kropf-Prodinger)

Average # of branches in a random plane tree of size n is

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 \triangleright δ ... periodic fluctuation:

$$\delta(x) := \frac{1}{\log 2} \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1 + \chi_k) \Gamma(\chi_k/2) \zeta(-1 + \chi_k) e^{2k\pi i x}, \quad \chi_k = \frac{2\pi i k}{\log 2}.$$

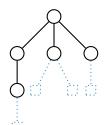


Leaves

$$\mathbb{E} \sim \frac{n}{r+1}$$

$$\mathbb{V} \sim \frac{r(r+2)}{6(r+1)^2}n$$

limit law: ✓

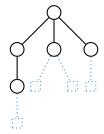




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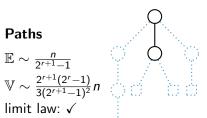
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Paths

$$\mathbb{E} \sim \frac{n}{2^{r+1}-1}$$

$$2^{r+1}(2^r-1)$$



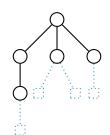


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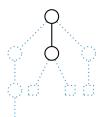


Paths

$$\mathbb{E} \sim \frac{n}{2^{r+1}-1}$$

$$\mathbb{V} \sim \frac{2^{r+1}(2^r-1)}{3(2^{r+1}-1)^2}n$$

limit law: ✓

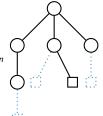


Old leaves

$$\mathbb{E} \sim (2 - B_{r-1}(1/4))n$$

$$\mathbb{V} = \Theta(n)$$

limit law: √



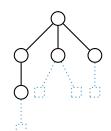


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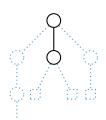


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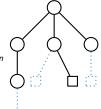


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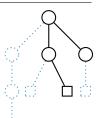


Old paths

$$\mathbb{E} \sim \frac{2n}{r+2}$$

$$\mathbb{V} \sim \frac{2r(r+1)}{3(r+2)^2}n$$

limit law: ✓

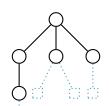


Leaves

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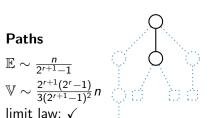
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Paths

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Disclaimer

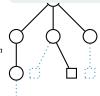
Results are not always that nice!

Old leaves

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$$\mathbb{V} = \Theta(n)$$

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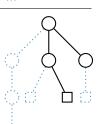


Old paths

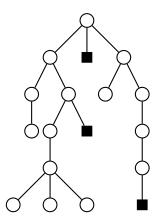
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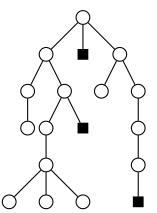


- ► Motivation: Stanley's Catalan interpretation #26
- ▶ Rightmost leaves in all branches of root have odd distance



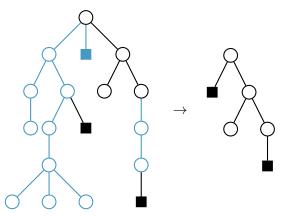


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- ▶ Reduction: remove parent & grandparent (except root) of ■





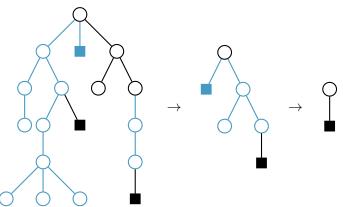
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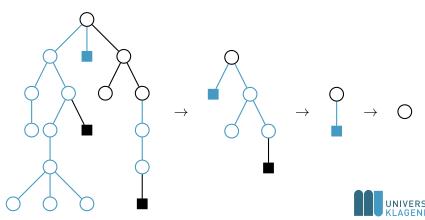


Further Models

- Motivation: Stanley's Catalan interpretation #26
- Rightmost leaves in all branches of root have odd distance
- Reduction: remove parent & grandparent (except root) of ■



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Counterexample: Results

► Reduction with different parameter behavior ✓

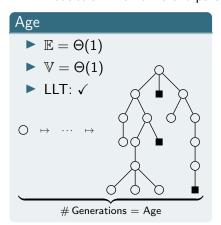
Age

Size of rth Reduction



Counterexample: Results

► Reduction with different parameter behavior ✓

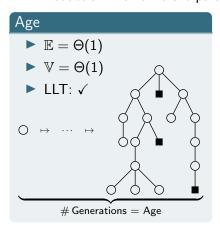


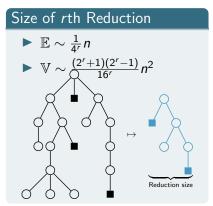
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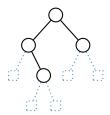
Cutting strategy:

- ▶ Remove Leaves
- ▶ Merge single children with their corresponding parent



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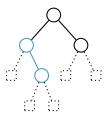




Further Models

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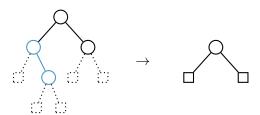




Further Models

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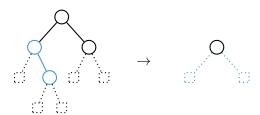
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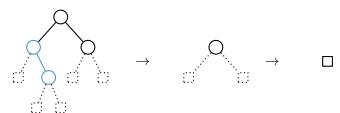




Further Models

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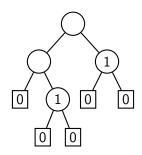
- ightharpoonup Leaves ightharpoonup 0
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- Otherwise: maximum of children



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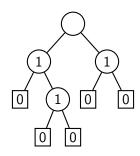


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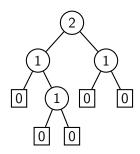
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We label the nodes according to the following rules:

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Further Models



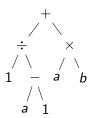
Age → Register function (Horton-Strahler-Index)

► Applications:



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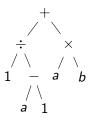
- ► Applications:
 - ▶ Required stack size for evaluating arithmetic expressions





Further Models

- ► Applications:
 - ▶ Required stack size for evaluating arithmetic expressions
 - ▶ Branching complexity of river networks (e.g. Danube: 9)

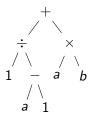






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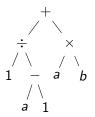




Asymptotic analysis:



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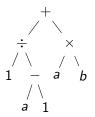




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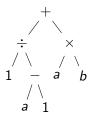




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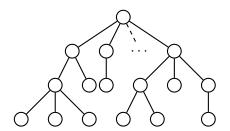


- Asymptotic analysis:
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 - r-branches, Numerics: Yamamoto, Yamazaki (2009)



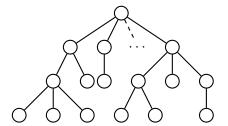
Core Idea: # removed vertices . . . additive tree parameter

 $ightharpoonup au \in \mathcal{T}$ tree; au_1, au_2, \ldots, au_k branches of au



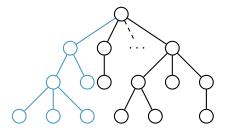


- $ightharpoonup au \in \mathcal{T}$ tree; au_1, au_2, \ldots, au_k branches of au
- $F(\tau) = F(\tau_1) + F(\tau_2) + \cdots + F(\tau_k) + f(\tau),$
 - \blacktriangleright with toll function $f: \mathcal{T} \to \mathbb{R}$



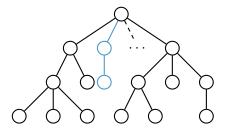


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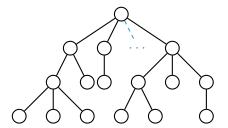


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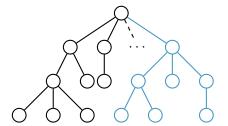


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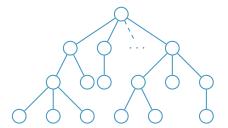


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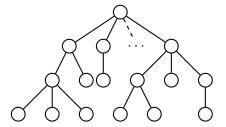


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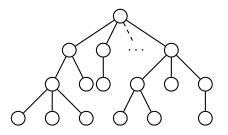


- ▶ Wagner (2015), Janson (2016), Wagner et al. (2018)...:
 - $ightharpoonup au_n$ random tree, size n; f suitable $\rightsquigarrow F(\tau_n)$ asymptotically Gaussian



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Families suitable for this approach:

- simply generated
- Pólya
- non-crossing
- **▶** ... (?)
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