

# Cutting Down and Growing Trees

Joint work with  
*Clemens Heuberger, Sara Kropf, Helmut Prodingler, Stephan Wagner*

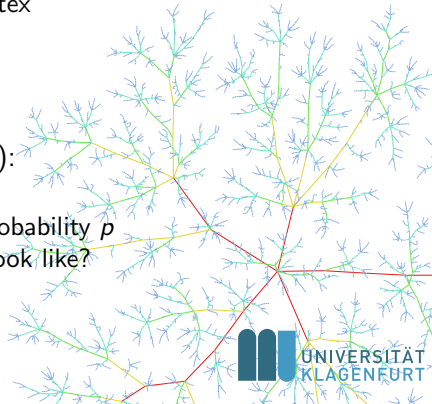


This work is licensed under a Creative Commons Attribution-NonCommercial-ShareAlike 4.0 International License.

# Related Research



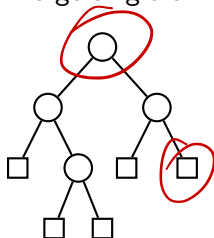
- ▶ Random edge removal (Meir–Moon, '70; Panholzer '06; ...)
  - ▶ Choose and remove random edge
  - ▶ Keep component with root vertex
  - ▶ How long does tree survive?
- ▶ Tree percolation (Lyons, '90, ...):
  - ▶ Infinite tree branching process
  - ▶ Remove edges (indep.) with probability  $p$
  - ▶ How does “root component” look like?



# Example: Trimming Binary Trees

Cutting strategy:

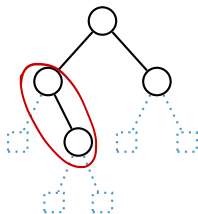
- ▶ Remove Leaves
- ▶ Merge single children with their corresponding parent



# Example: Trimming Binary Trees

Cutting strategy:

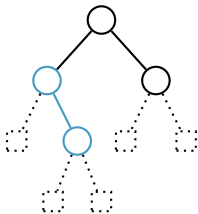
- ▶ Remove Leaves
- ▶ Merge single children with their corresponding parent



## Example: Trimming Binary Trees

Cutting strategy:

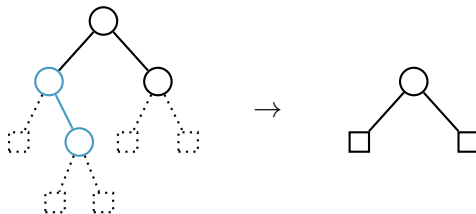
- ▶ Remove Leaves
- ▶ Merge single children with their corresponding parent



# Example: Trimming Binary Trees

Cutting strategy:

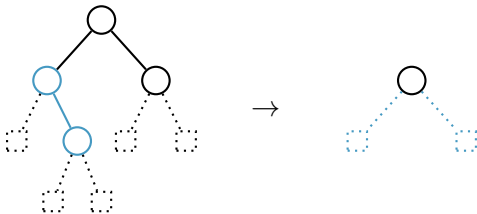
- ▶ Remove Leaves
- ▶ Merge single children with their corresponding parent



## Example: Trimming Binary Trees

Cutting strategy:

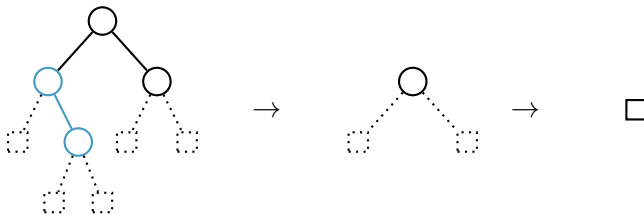
- ▶ Remove Leaves
- ▶ Merge single children with their corresponding parent



# Example: Trimming Binary Trees

Cutting strategy:

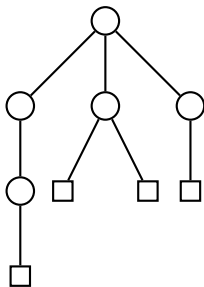
- ▶ Remove Leaves
- ▶ Merge single children with their corresponding parent





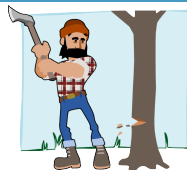
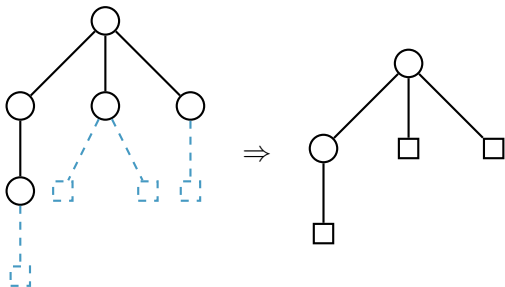
## Example: Cutting Down Plane Trees

- Remove all leaves!



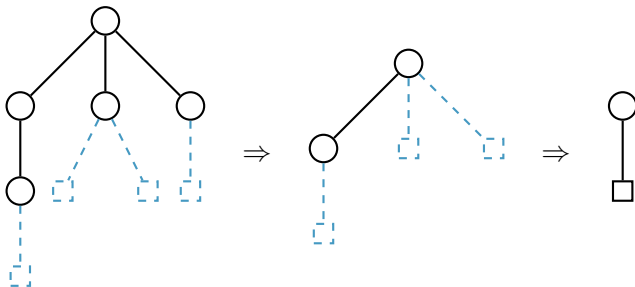
# Example: Cutting Down Plane Trees

- Remove all leaves!



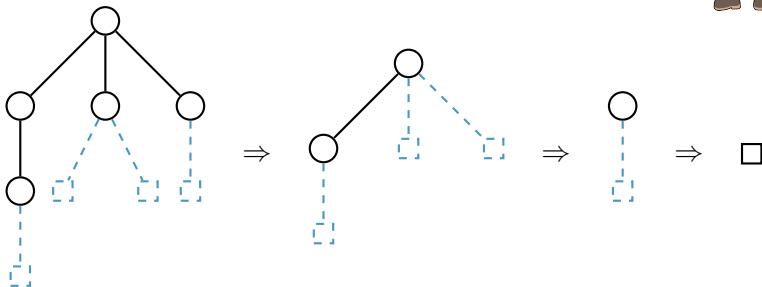
# Example: Cutting Down Plane Trees

- Remove all leaves!



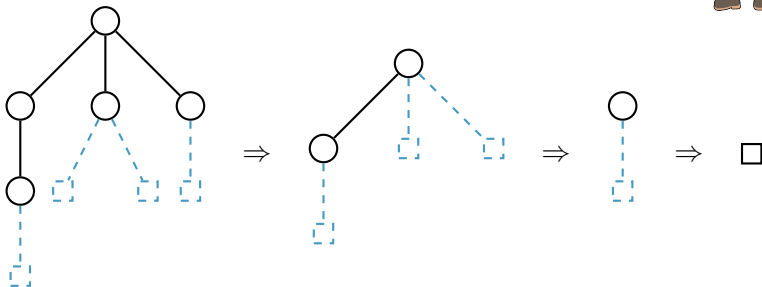
# Example: Cutting Down Plane Trees

- Remove all leaves!



# Example: Cutting Down Plane Trees

- Remove all leaves!



## Parameters of Interest:

- Size of  $r$ th reduction
- Age: # of possible reductions

## Reduction → Expansion

- ▶ modelling reduction directly: not suitable
- ▶ instead: see inverse operation, **growing trees**



.....

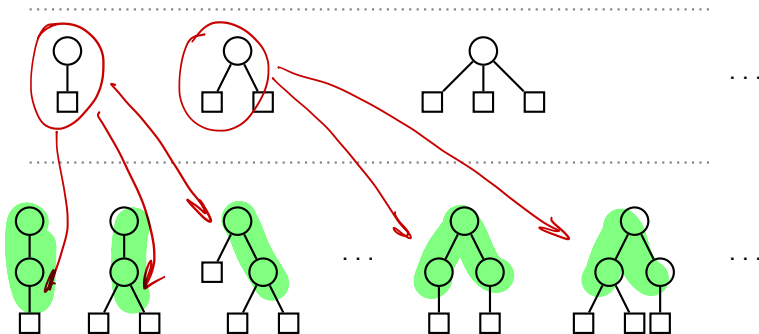
## Reduction → Expansion

- ▶ modelling reduction directly: not suitable
- ▶ instead: see inverse operation, **growing trees**



# Reduction → Expansion

- ▶ modelling reduction directly: not suitable
- ▶ instead: see inverse operation, **growing trees**





# Expansion operators

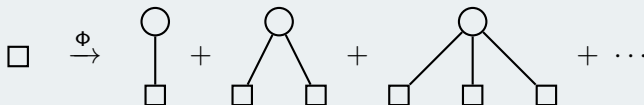
- ▶  $F \dots$  family of plane trees; bivariate generating function  $f$
- ▶ expansion operator  $\Phi \Rightarrow \Phi(f)$  counts expanded trees

# Expansion operators

- ▶  $F \dots$  family of plane trees; bivariate generating function  $f$
- ▶ expansion operator  $\Phi \Rightarrow \Phi(f)$  counts expanded trees

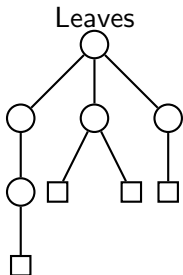
## Leaf expansion $\Phi$

- ▶ inverse operation to leaf reduction
  - ▶ attach leaves to all current leaves (**required**)
  - ▶ attach leaves to inner nodes (**optional**)

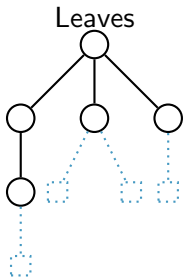


$$\square \triangleq t, \quad \bigcirc \triangleq z \quad \Rightarrow \quad \Phi(t) = zt + zt^2 + zt^3 + \dots$$

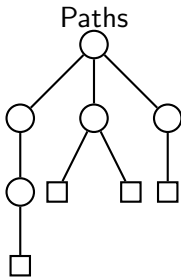
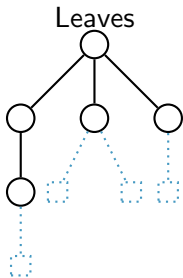
## Reductions on Plane Trees



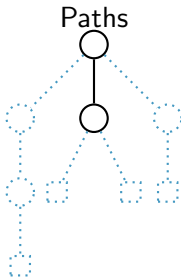
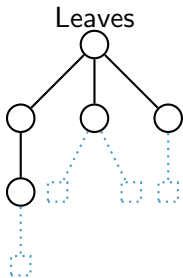
# Reductions on Plane Trees



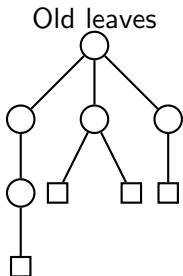
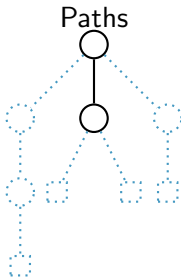
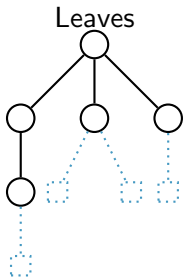
## Reductions on Plane Trees



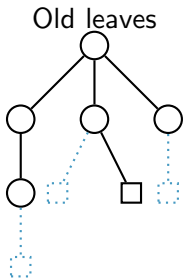
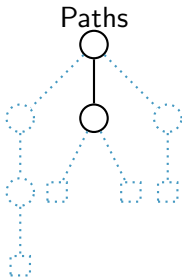
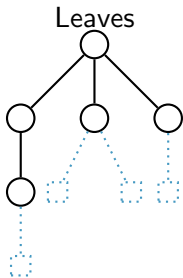
## Reductions on Plane Trees



## Reductions on Plane Trees



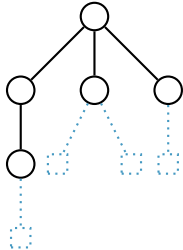
## Reductions on Plane Trees



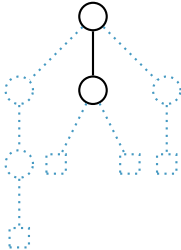


# Reductions on Plane Trees

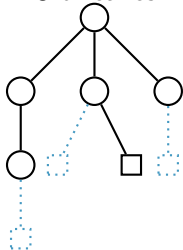
Leaves



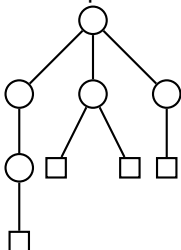
Paths



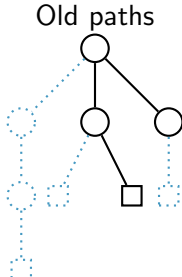
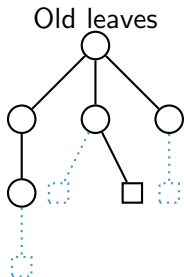
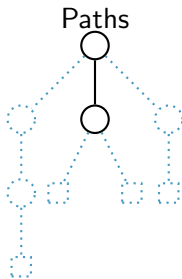
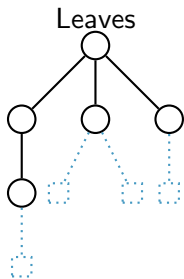
Old leaves



Old paths

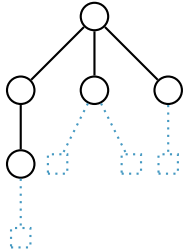


## Reductions on Plane Trees

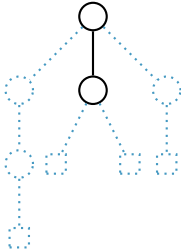


# Reductions on Plane Trees

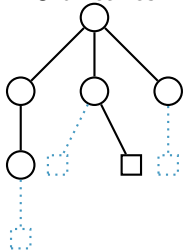
Leaves



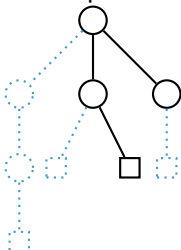
Paths



Old leaves



Old paths



## Parameters of Interest:

- ▶ tree size after  $r$  reductions
- ▶ cumulative reduction size

# Bivariate Generating Function

## Proposition

- ▶  $\mathcal{T} \dots$  rooted plane trees

# Bivariate Generating Function

## Proposition

- ▶  $\mathcal{T} \dots$  rooted plane trees
- ▶  $T(z, t) \dots$  BGF for  $\mathcal{T}$  ( $z \rightsquigarrow$  inner nodes,  $t \rightsquigarrow$  leaves)

# Bivariate Generating Function

## Proposition

- ▶  $\mathcal{T} \dots$  rooted plane trees
- ▶  $T(z, t) \dots$  BGF for  $\mathcal{T}$  ( $z \rightsquigarrow$  inner nodes,  $t \rightsquigarrow$  leaves)

$$\Rightarrow T(z, t) = \frac{1 - (z - t) - \sqrt{1 - 2(z + t) + (z - t)^2}}{2}$$

**Proof.** Symbolic equation

$$\mathcal{T} = \square + \begin{array}{c} \bigcirc \\ \swarrow \quad \downarrow \quad \searrow \\ \mathcal{T} \quad \mathcal{T} \quad \dots \quad \mathcal{T} \end{array}$$

translates into

$$T(z, t) = t + z \cdot \frac{T(z, t)}{1 - T(z, t)}$$

which can be solved explicitly.

# Bivariate Generating Function

## Proposition

- ▶  $\mathcal{T} \dots$  rooted plane trees
- ▶  $T(z, t) \dots$  BGF for  $\mathcal{T}$  ( $z \rightsquigarrow$  inner nodes,  $t \rightsquigarrow$  leaves)

$$\Rightarrow T(z, t) = \frac{1 - (z - t) - \sqrt{1 - 2(z + t) + (z - t)^2}}{2}$$

**Proof.** Symbolic equation

$$\mathcal{T} = \square + \begin{array}{c} \bigcirc \\ \swarrow \quad \downarrow \quad \searrow \\ \mathcal{T} \quad \mathcal{T} \quad \dots \quad \mathcal{T} \end{array}$$

translates into

$$T(z, t) = t + z \cdot \frac{T(z, t)}{1 - T(z, t)}$$

which can be solved explicitly.

# Bivariate Generating Function

## Proposition

- ▶  $\mathcal{T} \dots$  rooted plane trees
- ▶  $T(z, t) \dots$  BGF for  $\mathcal{T}$  ( $z \rightsquigarrow$  inner nodes,  $t \rightsquigarrow$  leaves)

$$\Rightarrow T(z, t) = \frac{1 - (z - t) - \sqrt{1 - 2(z + t) + (z - t)^2}}{2}$$

**Proof.** Symbolic equation

$$\mathcal{T} = \square + \begin{array}{c} \bigcirc \\ \swarrow \quad \downarrow \quad \searrow \\ \mathcal{T} \quad \mathcal{T} \quad \dots \quad \mathcal{T} \end{array}$$

translates into

$$T(z, t) = t + z \cdot \frac{T(z, t)}{1 - T(z, t)}$$

which can be solved explicitly.



# Bivariate Generating Function

## Proposition

- ▶  $\mathcal{T} \dots$  rooted plane trees
- ▶  $T(z, t) \dots$  BGF for  $\mathcal{T}$  ( $z \rightsquigarrow$  inner nodes,  $t \rightsquigarrow$  leaves)

$$\Rightarrow T(z, t) = \frac{1 - (z - t) - \sqrt{1 - 2(z + t) + (z - t)^2}}{2}$$

**Proof.** Symbolic equation

$$\mathcal{T} = \square + \begin{array}{c} \text{red circle} \\ \swarrow \quad \downarrow \quad \searrow \\ \mathcal{T} \quad \mathcal{T} \quad \dots \quad \mathcal{T} \end{array}$$

translates into

$$T(z, t) = t + z \cdot \frac{T(z, t)}{1 - T(z, t)}$$

which can be solved explicitly.

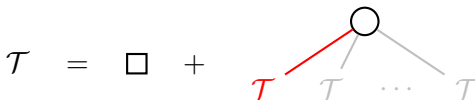
# Bivariate Generating Function

## Proposition

- ▶  $\mathcal{T} \dots$  rooted plane trees
- ▶  $T(z, t) \dots$  BGF for  $\mathcal{T}$  ( $z \rightsquigarrow$  inner nodes,  $t \rightsquigarrow$  leaves)

$$\Rightarrow T(z, t) = \frac{1 - (z - t) - \sqrt{1 - 2(z + t) + (z - t)^2}}{2}$$

**Proof.** Symbolic equation



translates into

$$T(z, t) = t + z \cdot \frac{T(z, t)}{1 - T(z, t)}$$

which can be solved explicitly.

# Bivariate Generating Function

## Proposition

- ▶  $\mathcal{T} \dots$  rooted plane trees
- ▶  $T(z, t) \dots$  BGF for  $\mathcal{T}$  ( $z \rightsquigarrow$  inner nodes,  $t \rightsquigarrow$  leaves)

$$\Rightarrow T(z, t) = \frac{1 - (z - t) - \sqrt{1 - 2(z + t) + (z - t)^2}}{2}$$

**Proof.** Symbolic equation

$$\mathcal{T} = \square + \begin{array}{c} \circ \\ \swarrow \quad \downarrow \quad \searrow \\ \mathcal{T} \quad \mathcal{T} \quad \cdots \quad \mathcal{T} \end{array}$$

translates into

$$T(z, t) = t + z \cdot \frac{T(z, t)}{1 - T(z, t)}$$

which can be solved explicitly.

# Bivariate Generating Function

## Proposition

- ▶  $\mathcal{T} \dots$  rooted plane trees
- ▶  $T(z, t) \dots$  BGF for  $\mathcal{T}$  ( $z \rightsquigarrow$  inner nodes,  $t \rightsquigarrow$  leaves)

$$\Rightarrow T(z, t) = \frac{1 - (z - t) - \sqrt{1 - 2(z + t) + (z - t)^2}}{2}$$

**Proof.** Symbolic equation

$$\mathcal{T} = \square + \begin{array}{c} \bigcirc \\ \swarrow \quad \downarrow \quad \searrow \\ \mathcal{T} \quad \mathcal{T} \quad \dots \quad \mathcal{T} \end{array}$$

translates into

$$T(z, t) = t + z \cdot \frac{T(z, t)}{1 - T(z, t)}$$

which can be solved explicitly.

# SageMath Demo

# Leaf expansion operator $\Phi$

## Proposition

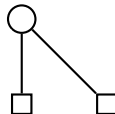
$$\Phi(f(z, t)) = (1 - t)f\left(\frac{z}{(1 - t)^2}, \frac{zt}{(1 - t)^2}\right)$$

# Leaf expansion operator $\Phi$

## Proposition

$$\Phi(f(z, t)) = (1 - t)f\left(\frac{z}{(1 - t)^2}, \frac{zt}{(1 - t)^2}\right)$$

- ▶ Tree with  $n$  inner nodes and  $k$  leaves  $\rightsquigarrow z^n t^k$
- ▶ **Expansion:**



- ▶ In total:

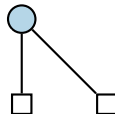
$$\Phi(z^n t^k) =$$

# Leaf expansion operator $\Phi$

## Proposition

$$\Phi(f(z, t)) = (1 - t)f\left(\frac{z}{(1 - t)^2}, \frac{zt}{(1 - t)^2}\right)$$

- ▶ Tree with  $n$  inner nodes and  $k$  leaves  $\rightsquigarrow z^n t^k$
- ▶ **Expansion:**
  - ▶ inner nodes **stay** inner nodes



- ▶ In total:

$$\Phi(z^n t^k) = z^n.$$



# Leaf expansion operator $\Phi$

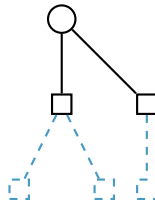
## Proposition

$$\Phi(f(z, t)) = (1 - t)f\left(\frac{z}{(1 - t)^2}, \frac{zt}{(1 - t)^2}\right)$$

- ▶ Tree with  $n$  inner nodes and  $k$  leaves  $\rightsquigarrow z^n t^k$

- ▶ **Expansion:**

- ▶ inner nodes **stay** inner nodes
- ▶ attach a **non-empty sequence of leaves** to all current leaves



- ▶ In total:

$$\Phi(z^n t^k) = z^n \cdot \left(\frac{zt}{1 - t}\right)^k.$$

# Leaf expansion operator $\Phi$

## Proposition

$$\Phi(f(z, t)) = (1 - t)f\left(\frac{z}{(1 - t)^2}, \frac{zt}{(1 - t)^2}\right)$$

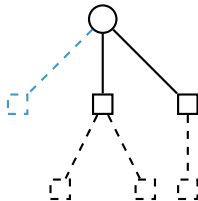
- ▶ Tree with  $n$  inner nodes and  $k$  leaves  $\rightsquigarrow z^n t^k$

- ▶ **Expansion:**

- ▶ inner nodes **stay** inner nodes
- ▶ attach a **non-empty sequence of leaves** to all current leaves
- ▶ there are  $2n + k - 1$  positions where sequences of leaves can be inserted

- ▶ In total:

$$\Phi(z^n t^k) = z^n \cdot \left(\frac{zt}{1 - t}\right)^k \cdot \frac{1}{(1 - t)^{2n+k-1}}$$



# Leaf expansion operator $\Phi$

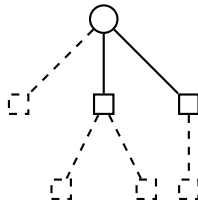
## Proposition

$$\Phi(f(z, t)) = (1 - t)f\left(\frac{z}{(1 - t)^2}, \frac{zt}{(1 - t)^2}\right)$$

- ▶ Tree with  $n$  inner nodes and  $k$  leaves  $\rightsquigarrow z^n t^k$

- ▶ **Expansion:**

- ▶ inner nodes **stay** inner nodes
- ▶ attach a **non-empty sequence of leaves** to all current leaves
- ▶ there are  $2n + k - 1$  positions where sequences of leaves can be inserted



- ▶ In total:

$$\Phi(z^n t^k) = z^n \cdot \left(\frac{zt}{1 - t}\right)^k \cdot \frac{1}{(1 - t)^{2n + k - 1}} = (1 - t) \left(\frac{z}{(1 - t)^2}\right)^n \left(\frac{zt}{(1 - t)^2}\right)^k$$

# Leaf expansion operator $\Phi$

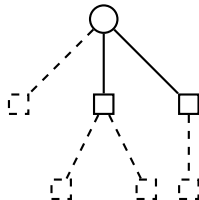
## Proposition

$$\Phi(f(z, t)) = (1 - t)f\left(\frac{z}{(1 - t)^2}, \frac{zt}{(1 - t)^2}\right)$$

- ▶ Tree with  $n$  inner nodes and  $k$  leaves  $\rightsquigarrow z^n t^k$

- ▶ **Expansion:**

- ▶ inner nodes **stay** inner nodes
- ▶ attach a **non-empty sequence of leaves** to all current leaves
- ▶ there are  $2n + k - 1$  positions where sequences of leaves can be inserted



- ▶ In total:

$$\Phi(z^n t^k) = z^n \cdot \left(\frac{zt}{1 - t}\right)^k \cdot \frac{1}{(1 - t)^{2n + k - 1}} = (1 - t) \left(\frac{z}{(1 - t)^2}\right)^n \left(\frac{zt}{(1 - t)^2}\right)^k$$

- ▶ As  $\Phi$  is linear, this proves the proposition.

# Properties of $\Phi$

- ▶ Functional equation:  $T(z, t) = \Phi(T(z, t)) + t$

# Properties of $\Phi$

- ▶ Functional equation:  $T(z, t) = \Phi(T(z, t)) + t$
- ▶ With  $z = u/(1 + u)^2$  and by some manipulations

$$\Phi^r(T(z, t))|_{t=z} = \frac{1 - u^{r+2}}{(1 - u^{r+1})(1 + u)} T\left(\frac{u(1 - u^{r+1})^2}{(1 - u^{r+2})^2}, \frac{u^{r+1}(1 - u)^2}{(1 - u^{r+2})^2}\right)$$

# Properties of $\Phi$

- ▶ Functional equation:  $T(z, t) = \Phi(T(z, t)) + t$
- ▶ With  $z = u/(1 + u)^2$  and by some manipulations

$$\Phi^r(T(z, t))|_{t=z} = \frac{1 - u^{r+2}}{(1 - u^{r+1})(1 + u)} T\left(\frac{u(1 - u^{r+1})^2}{(1 - u^{r+2})^2}, \frac{u^{r+1}(1 - u)^2}{(1 - u^{r+2})^2}\right)$$

- ▶ BGF  $G_r(z, v)$  for size comparison:  $z$  tracks original size,  $v$  size of  $r$ -fold reduced tree

## Properties of $\Phi$

- ▶ Functional equation:  $T(z, t) = \Phi(T(z, t)) + t$
- ▶ With  $z = u/(1 + u)^2$  and by some manipulations

$$\Phi^r(T(z, t))|_{t=z} = \frac{1 - u^{r+2}}{(1 - u^{r+1})(1 + u)} T\left(\frac{u(1 - u^{r+1})^2}{(1 - u^{r+2})^2}, \frac{u^{r+1}(1 - u)^2}{(1 - u^{r+2})^2}\right)$$

- ▶ BGF  $G_r(z, v)$  for size comparison:  $z$  tracks original size,  $v$  size of  $r$ -fold reduced tree
- ▶ Intuition:  $v$  “remembers” size while tree family is expanded

$$G_r(z, v) = \frac{1 - u^{r+2}}{(1 - u^{r+1})(1 + u)} T\left(\frac{u(1 - u^{r+1})^2}{(1 - u^{r+2})^2} v, \frac{u^{r+1}(1 - u)^2}{(1 - u^{r+2})^2} v\right)$$



# Cutting leaves

## Theorem (H.–Heuberger–Kropf–Prodinger)

*After  $r$  reductions of a random tree of size  $n$ , the remaining size  $X_{n,r}$  has **mean** and **variance***

$$\mathbb{E}X_{n,r} = \frac{n}{r+1} - \frac{r(r-1)}{6(r+1)} + O(n^{-1}),$$

$$\mathbb{V}X_{n,r} = \frac{r(r+2)}{6(r+1)^2}n + O(1),$$

*and  $X_{n,r}$  is **asymptotically normally distributed**.*

# Cutting leaves

## Theorem (H.–Heuberger–Kropf–Prodinger)

After  $r$  reductions of a random tree of size  $n$ , the remaining size  $X_{n,r}$  has **mean** and **variance**

$$\mathbb{E}X_{n,r} = \frac{n}{r+1} - \frac{r(r-1)}{6(r+1)} + O(n^{-1}),$$

$$\mathbb{V}X_{n,r} = \frac{r(r+2)}{6(r+1)^2}n + O(1),$$

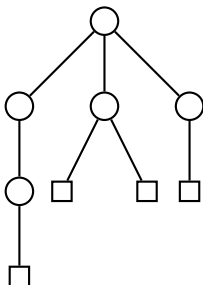
and  $X_{n,r}$  is **asymptotically normally distributed**.

### Proof insights:

- ▶  $\mathbb{E}X_{n,r}$  and  $\mathbb{V}X_{n,r}$  follow via singularity analysis
- ▶ Asymptotic normality:  $n - X_{n,r}$  is a **tree parameter** with small **toll function**, limit law by Wagner (2015)

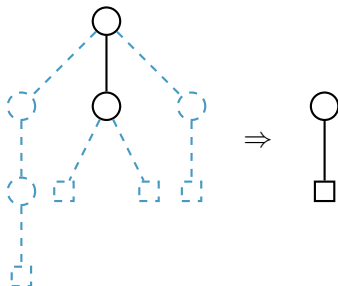
# Pruning

- Remove all paths that end in a leaf!



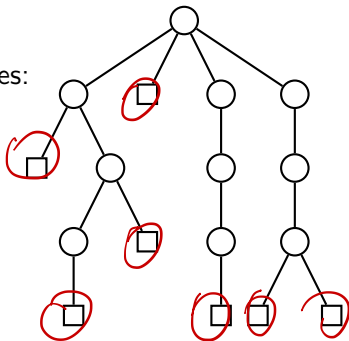
# Pruning

- Remove all paths that end in a leaf!



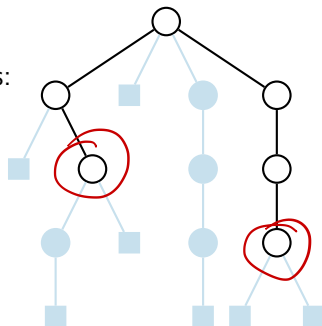
# Branches in a Tree

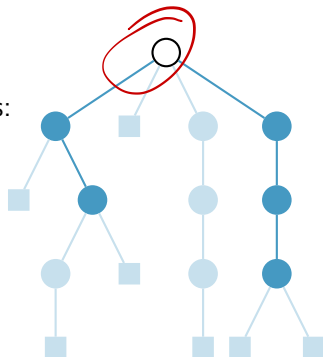
- Trees can be partitioned into branches:

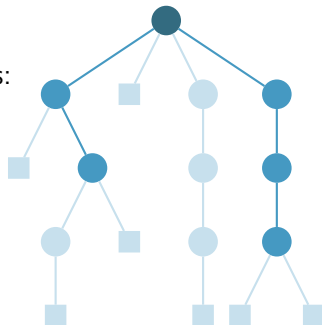


# Branches in a Tree

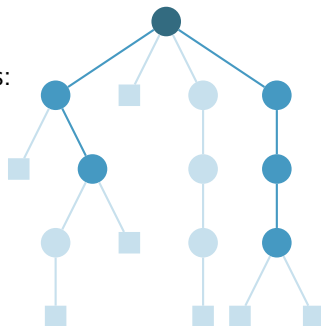
- Trees can be partitioned into branches:

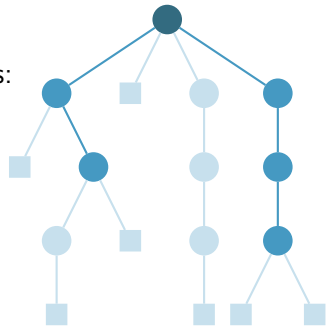


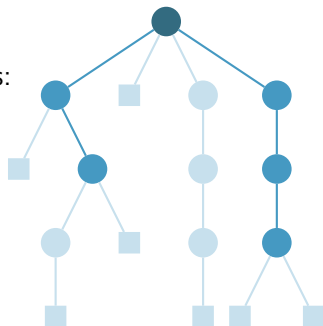










☐

# Branches in a Tree – Result

# Branches in a Tree – Result

## Theorem (H.–Heuberger–Kropf–Prodinger)

*Average # of branches in a random plane tree of size  $n$  is*

$$\alpha n + \frac{1}{6} \log_4 n + C + \delta(\log_4 n) + O(n^{-1/4}),$$

# Branches in a Tree – Result

## Theorem (H.–Heuberger–Kropf–Prodinger)

*Average # of branches in a random plane tree of size  $n$  is*

$$\alpha n + \frac{1}{6} \log_4 n + C + \delta(\log_4 n) + O(n^{-1/4}),$$

►  $\alpha = \sum_{k \geq 2} \frac{1}{2^k - 1} \approx 0.60669,$

# Branches in a Tree – Result

## Theorem (H.–Heuberger–Kropf–Prodinger)

*Average # of branches in a random plane tree of size  $n$  is*

$$\alpha n + \frac{1}{6} \log_4 n + C + \delta(\log_4 n) + O(n^{-1/4}),$$

$$\blacktriangleright \alpha = \sum_{k \geq 2} \frac{1}{2^k - 1} \approx 0.60669,$$

$$\blacktriangleright C = -\frac{\gamma + 4\alpha \log 2 + \log 2 + 24\zeta'(-1) + 2}{12 \log 2} \approx -0.11811,$$

# Branches in a Tree – Result

## Theorem (H.–Heuberger–Kropf–Prodinger)

*Average # of branches in a random plane tree of size  $n$  is*

$$\alpha n + \frac{1}{6} \log_4 n + C + \delta(\log_4 n) + O(n^{-1/4}),$$

▶  $\alpha = \sum_{k \geq 2} \frac{1}{2^k - 1} \approx 0.60669,$

▶  $C = -\frac{\gamma + 4\alpha \log 2 + \log 2 + 24\zeta'(-1) + 2}{12 \log 2} \approx -0.11811,$

▶  $\delta \dots$  periodic fluctuation:

$$\delta(x) := \frac{1}{\log 2} \sum_{k \in \mathbb{Z} \setminus \{0\}} (-1 + \chi_k) \Gamma(\chi_k/2) \zeta(-1 + \chi_k) e^{2k\pi ix}, \quad \chi_k = \frac{2\pi ik}{\log 2}.$$



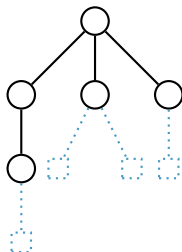
# Summary: Reductions on Plane Trees

## Leaves

$$\mathbb{E} \sim \frac{n}{r+1}$$

$$\mathbb{V} \sim \frac{r(r+2)}{6(r+1)^2} n$$

limit law: ✓



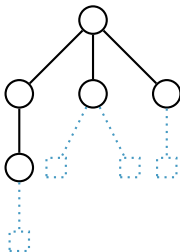
# Summary: Reductions on Plane Trees

## Leaves

$$\mathbb{E} \sim \frac{n}{r+1}$$

$$\mathbb{V} \sim \frac{r(r+2)}{6(r+1)^2} n$$

limit law: ✓

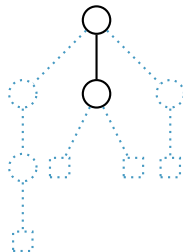


## Paths

$$\mathbb{E} \sim \frac{n}{2^{r+1}-1}$$

$$\mathbb{V} \sim \frac{2^{r+1}(2^r-1)}{3(2^{r+1}-1)^2} n$$

limit law: ✓



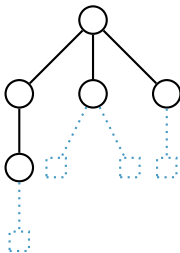
## Summary: Reductions on Plane Trees

## Leaves

$$\mathbb{E} \sim \frac{n}{r+1}$$

$$\mathbb{V} \sim \frac{r(r+2)}{6(r+1)^2} n$$

limit law: ✓

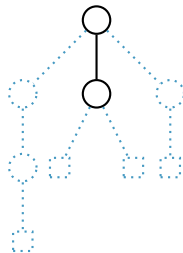


## Paths

$$\mathbb{E} \sim \frac{n}{2^{r+1}-1}$$

$$\mathbb{V} \sim \frac{2^{r+1}(2^r-1)}{3(2^{r+1}-1)^2} n$$

limit law: ✓

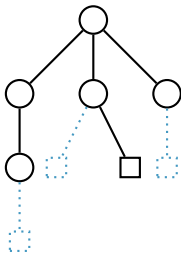


## Old leaves

$$\mathbb{E} \sim (2 - B_{r-1}(1/4))n$$

$$\mathbb{V} = \Theta(n)$$

limit law: ✓



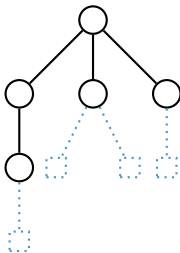
## Summary: Reductions on Plane Trees

## Leaves

$$\mathbb{E} \sim \frac{n}{r+1}$$

$$\mathbb{V} \sim \frac{r(r+2)}{6(r+1)^2} n$$

limit law: ✓

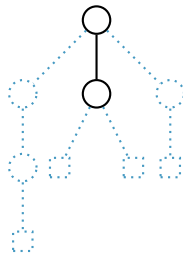


## Paths

$$\mathbb{E} \sim \frac{n}{2^{r+1}-1}$$

$$\mathbb{V} \sim \frac{2^{r+1}(2^r-1)}{3(2^{r+1}-1)^2} n$$

limit law: ✓

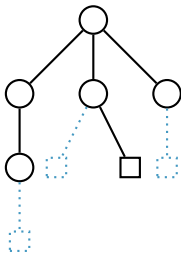


## Old leaves

$$\mathbb{E} \sim (2 - B_{r-1}(1/4))n$$

$$\mathbb{V} = \Theta(n)$$

limit law: ✓

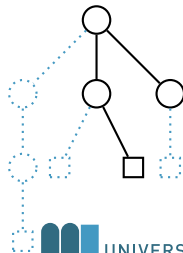


## Old paths

$$\mathbb{E} \sim \frac{2n}{r+2}$$

$$\mathbb{V} \sim \frac{2r(r+1)}{3(r+2)^2} n$$

limit law: ✓



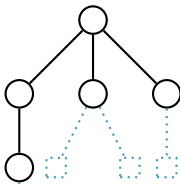
## Summary: Reductions on Plane Trees

## Leaves

$$\mathbb{E} \sim \frac{n}{r+1}$$

$$\mathbb{V} \sim \frac{r(r+2)}{6(r+1)^2} n$$

limit law: ✓

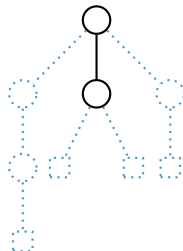


## Paths

$$\mathbb{E} \sim \frac{n}{2^{r+1}-1}$$

$$\mathbb{V} \sim \frac{2^{r+1}(2^r-1)}{3(2^{r+1}-1)^2} n$$

limit law: ✓



## Disclaimer

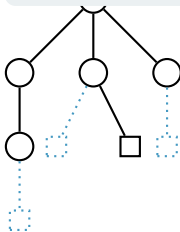
Results are **not always** that nice!

## Old leaves

$$\mathbb{E} \sim (2 - B_{r-1}(1/4))n$$

$$\mathbb{V} = \Theta(n)$$

limit law: ✓

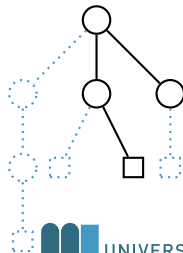


## Old paths

$$\mathbb{E} \sim \frac{2n}{r+2}$$

$$\mathbb{V} \sim \frac{2r(r+1)}{3(r+2)^2} n$$

limit law: ✓

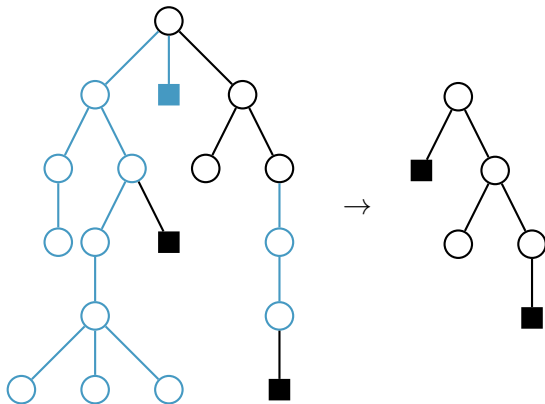






## Counterexample: Catalan–Stanley trees

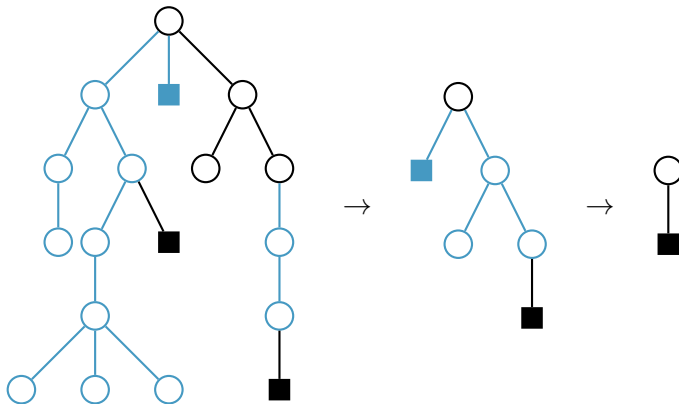
- ▶ Motivation: Stanley's Catalan interpretation #26
- ▶ Rightmost leaves in all branches of root have odd distance
- ▶ **Reduction:** remove parent/grandparent (except root) of ■





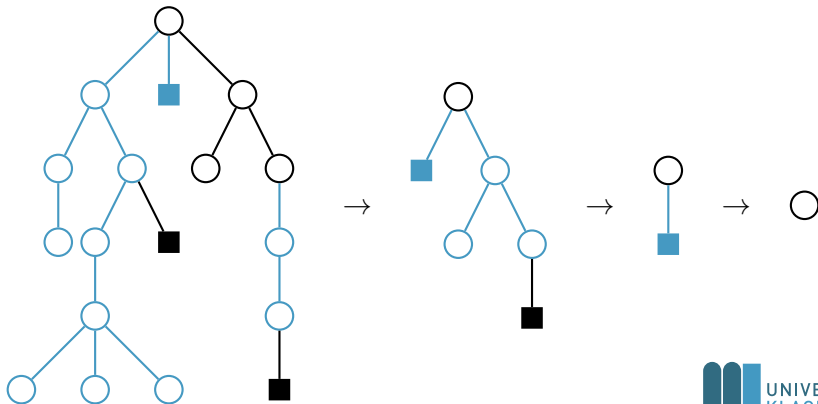
## Counterexample: Catalan–Stanley trees

- ▶ Motivation: Stanley's Catalan interpretation #26
- ▶ Rightmost leaves in all branches of root have odd distance
- ▶ **Reduction:** remove parent/grandparent (except root) of ■



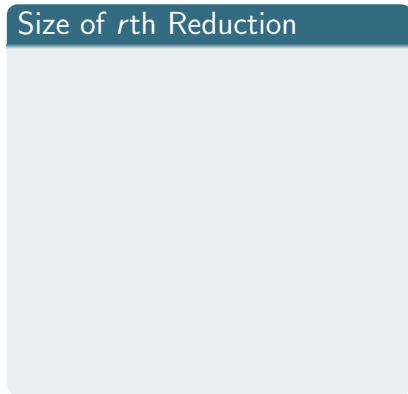
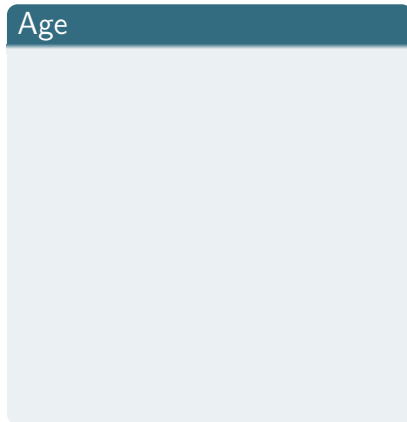
## Counterexample: Catalan–Stanley trees

- ▶ Motivation: Stanley's Catalan interpretation #26
- ▶ Rightmost leaves in all branches of root have odd distance
- ▶ **Reduction:** remove parent/grandparent (except root) of ■



## Counterexample: Results

- Reduction with different parameter behavior ✓

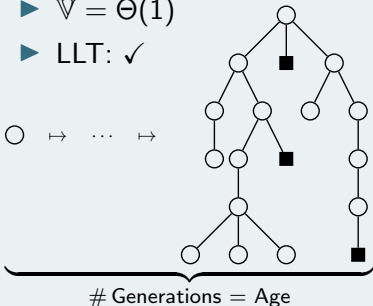


# Counterexample: Results

- ▶ Reduction with different parameter behavior ✓

## Age

- ▶  $\mathbb{E} = \Theta(1)$
- ▶  $\mathbb{V} = \Theta(1)$
- ▶ LLT: ✓



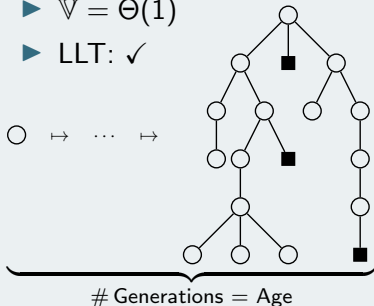
## Size of $r$ th Reduction

# Counterexample: Results

- ▶ Reduction with different parameter behavior ✓

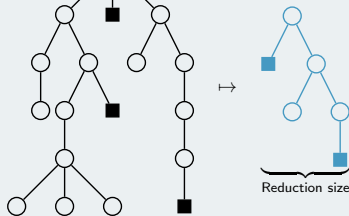
## Age

- ▶  $\mathbb{E} = \Theta(1)$
- ▶  $\mathbb{V} = \Theta(1)$
- ▶ LLT: ✓



## Size of $r$ th Reduction

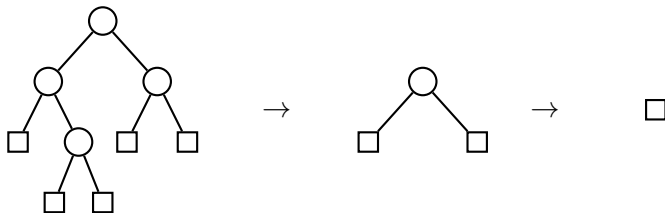
- ▶  $\mathbb{E} \sim \frac{1}{4^r} n$
- ▶  $\mathbb{V} \sim \frac{(2^r+1)(2^r-1)}{16^r} n^2$



# Trimming Binary Trees

Cutting strategy:

- ▶ Remove Leaves
- ▶ Merge single children with their corresponding parent



## Bonus: Touchard's Identity

### Proposition

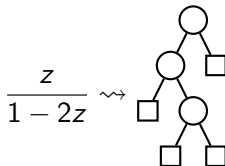
$$B(z) = 1 + \frac{z}{1-2z} B\left(\frac{z^2}{(1-2z)^2}\right)$$

## Bonus: Touchard's Identity

### Proposition

$$B(z) = 1 + \frac{z}{1-2z} B\left(\frac{z^2}{(1-2z)^2}\right)$$

- ▶ Binary trees consist of...
  - ▶ ...just a leaf,
  - ▶ ...a “smaller” tree with all leafs replaced by “chains”

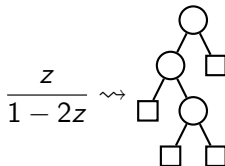




## Bonus: Touchard's Identity

### Proposition

$$B(z) = 1 + \frac{z}{1-2z} B\left(\frac{z^2}{(1-2z)^2}\right)$$



- ▶ Binary trees consist of...
  - ▶ ...just a leaf,
  - ▶ ...a “smaller” tree with all leaves replaced by “chains”

### Corollary (Touchard's Identity)

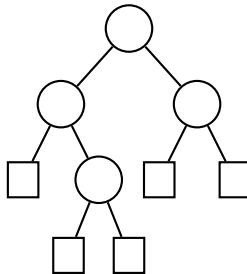
The Catalan numbers  $C_n = \frac{1}{n+1} \binom{2n}{n}$  satisfy

$$C_{n+1} = \sum_{0 \leq k \leq n/2} C_k 2^{n-2k} \binom{n}{2k}.$$

# How “old” do the nodes get?

We label the nodes according to the following rules:

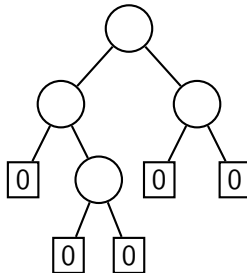
- ▶ Leaves  $\rightarrow 0$
- ▶  $\text{age}(\text{left child}) = \text{age}(\text{right child}) \rightarrow \text{increase by } 1$
- ▶ Otherwise: maximum of children



# How “old” do the nodes get?

We label the nodes according to the following rules:

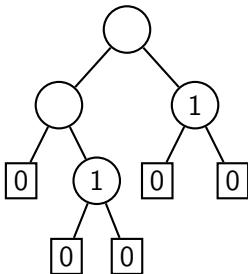
- ▶ Leaves  $\rightarrow 0$
- ▶  $\text{age}(\text{left child}) = \text{age}(\text{right child}) \rightarrow \text{increase by } 1$
- ▶ Otherwise: maximum of children



# How “old” do the nodes get?

We label the nodes according to the following rules:

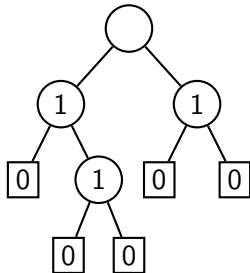
- ▶ Leaves  $\rightarrow 0$
- ▶  $\text{age}(\text{left child}) = \text{age}(\text{right child}) \rightarrow \text{increase by } 1$
- ▶ Otherwise: maximum of children



# How “old” do the nodes get?

We label the nodes according to the following rules:

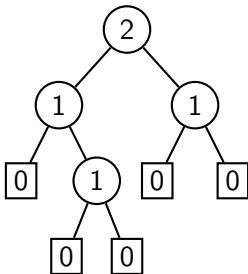
- ▶ Leaves  $\rightarrow 0$
- ▶  $\text{age}(\text{left child}) = \text{age}(\text{right child}) \rightarrow \text{increase by } 1$
- ▶ Otherwise: maximum of children



# How “old” do the nodes get?

We label the nodes according to the following rules:

- ▶ Leaves  $\rightarrow 0$
- ▶  $\text{age}(\text{left child}) = \text{age}(\text{right child}) \rightarrow \text{increase by } 1$
- ▶ Otherwise: maximum of children



# The Register Function

*Age  $\rightsquigarrow$  Register function (Horton-Strahler-Index)*

# The Register Function

Age  $\rightsquigarrow$  *Register function (Horton-Strahler-Index)*

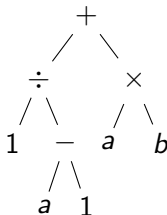
► Applications:



# The Register Function

Age  $\rightsquigarrow$  Register function (*Horton-Strahler-Index*)

- ▶ Applications:
  - ▶ Required stack size for evaluating arithmetic expressions

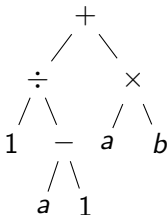


# The Register Function

Age  $\rightsquigarrow$  *Register function (Horton-Strahler-Index)*

► Applications:

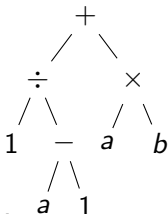
- Required stack size for evaluating arithmetic expressions
- Branching complexity of river networks (e.g. Danube: 9)



# The Register Function

Age  $\rightsquigarrow$  Register function (*Horton-Strahler-Index*)

- ▶ Applications:
  - ▶ Required stack size for evaluating arithmetic expressions
  - ▶ Branching complexity of river networks (e.g. Danube: 9)



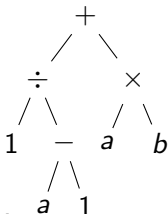
- ▶ Asymptotic analysis:



# The Register Function

Age  $\rightsquigarrow$  Register function (*Horton-Strahler-Index*)

- Applications:
  - Required stack size for evaluating arithmetic expressions
  - Branching complexity of river networks (e.g. Danube: 9)

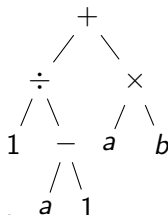


- Asymptotic analysis:
  - Flajolet, Raoult, Vuillemin (1979)

# The Register Function

Age  $\rightsquigarrow$  Register function (*Horton-Strahler-Index*)

- ▶ Applications:
  - ▶ Required stack size for evaluating arithmetic expressions
  - ▶ Branching complexity of river networks (e.g. Danube: 9)

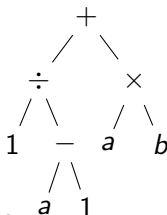


- ▶ Asymptotic analysis:
  - ▶ Flajolet, Raoult, Vuillemin (1979)
  - ▶ Flajolet, Prodinger (1986)

# The Register Function

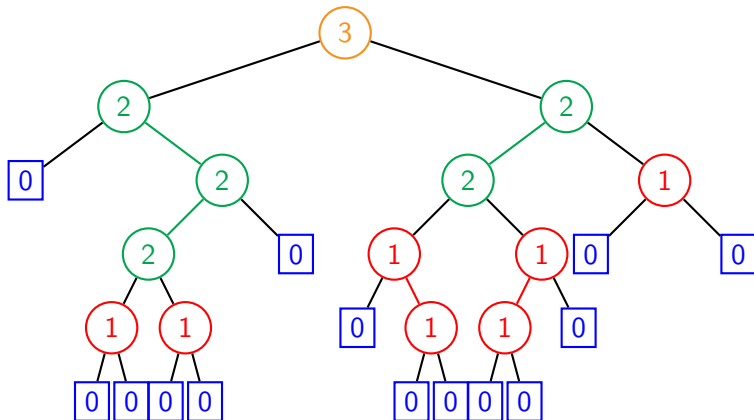
Age  $\rightsquigarrow$  Register function (*Horton-Strahler-Index*)

- ▶ Applications:
  - ▶ Required stack size for evaluating arithmetic expressions
  - ▶ Branching complexity of river networks (e.g. Danube: 9)

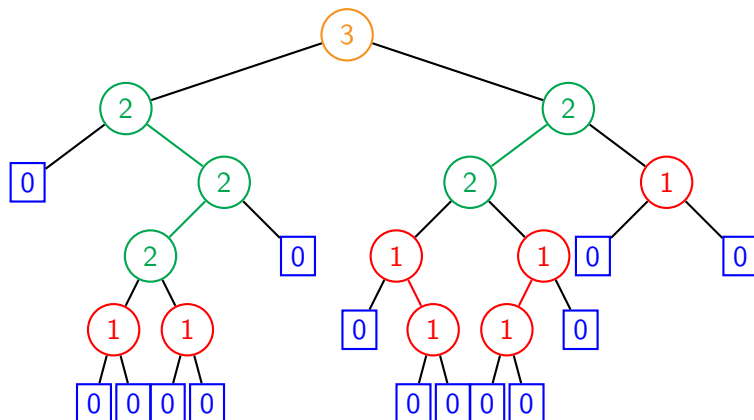


- ▶ Asymptotic analysis:
  - ▶ Flajolet, Raoult, Vuillemin (1979)
  - ▶ Flajolet, Prodinger (1986)
  - ▶  $r$ -branches, Numerics: Yamamoto, Yamazaki (2009)

# Local Structures – “ $r$ -branches”



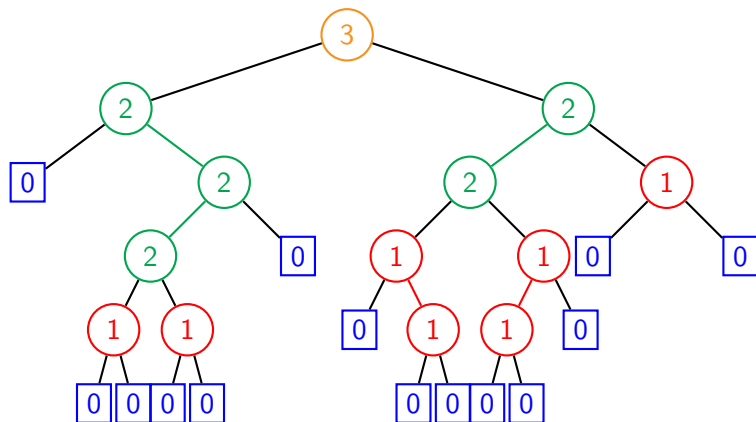
# Local Structures – “ $r$ -branches”



- Number / Distribution of ( $r$ -)branches?



## Local Structures – “ $r$ -branches”



- ▶ Number / Distribution of ( $r$ -)branches?

- Example:
- | $r$             | 0  | 1 | 2 | 3 |
|-----------------|----|---|---|---|
| # $r$ -branches | 14 | 5 | 2 | 1 |

## “ $r$ -branches” – Results

### Theorem (H.–Heuberger–Prodinger)

*In a random binary tree of size  $n$  . . .*

## “ $r$ -branches” – Results

### Theorem (H.–Heuberger–Prodinger)

*In a random binary tree of size  $n$  . . .*

- ▶ *# of  $r$ -branches is **asymptotically normally distributed***

## “ $r$ -branches” – Results

### Theorem (H.–Heuberger–Prodinger)

*In a random binary tree of size  $n$ ...*

- ▶ # of  $r$ -branches is **asymptotically normally distributed**
- ▶ with **mean** and **variance**

$$\mathbb{E} = \frac{n}{4^r} + \frac{1}{6} \left( 1 + \frac{5}{4^r} \right) + O(n^{-1}), \quad \mathbb{V} = \frac{4^r - 1}{3 \cdot 16^r} n + O(1)$$

# “ $r$ -branches” – Results

## Theorem (H.–Heuberger–Prodinger)

*In a random binary tree of size  $n$ ...*

- ▶ # of  $r$ -branches is **asymptotically normally distributed**
- ▶ with **mean** and **variance**

$$\mathbb{E} = \frac{n}{4^r} + \frac{1}{6} \left( 1 + \frac{5}{4^r} \right) + O(n^{-1}), \quad \mathbb{V} = \frac{4^r - 1}{3 \cdot 16^r} n + O(1)$$

- ▶ **expected total #** of branches is

$$\frac{4}{3}n + \frac{1}{6} \log_4 n + C + \delta(\log_4 n) + O(n^{-1} \log n),$$

- ▶  $C \approx 1.36190$ ,  $\delta \dots$  periodic fluctuation

# “ $r$ -branches” – Results

## Theorem (H.–Heuberger–Prodinger)

*In a random binary tree of size  $n$ ...*

- ▶ # of  $r$ -branches is **asymptotically normally distributed**
- ▶ with **mean** and **variance**

$$\mathbb{E} = \frac{n}{4^r} + \frac{1}{6} \left( 1 + \frac{5}{4^r} \right) + O(n^{-1}), \quad \mathbb{V} = \frac{4^r - 1}{3 \cdot 16^r} n + O(1)$$

- ▶ **expected total #** of branches is

$$\frac{4}{3}n + \frac{1}{6} \log_4 n + C + \delta(\log_4 n) + O(n^{-1} \log n),$$

- ▶  $C \approx 1.36190$ ,  $\delta \dots$  periodic fluctuation