## Benjamin HackI

## Cutting Down and Growing Trees

KARL
POPPER
KOLLEG

Der Wissenschaftsfonds.

## Is this about Gardening?

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- No. But Trees!


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"Asymptotic Analysis of Shape Parameters of Trees and Lattice Paths"


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## Analytic Combinatorics. . .

... deals with enumeration of discrete objects by using analytic methods.

$$
\frac{1-\sqrt{1-4 z}}{2}=1 z+1 z^{2}+2 z^{3}+5 z^{4}+14 z^{5}+42 z^{6}+\ldots
$$

## Trimming Binary Trees

Cutting strategy:

- Remove Leaves
- Merge single children with their corresponding parent



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We label the nodes according to the following rules:

- Leaves $\rightarrow 0$
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- Flajolet, Raoult, Vuillemin (1979)
- Flajolet, Prodinger (1986)
- r-Branches, Numerics: Yamamoto, Yamazaki (2009)

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\mathbb{E}=\frac{n}{4^{r}}+\frac{1}{6}\left(1+\frac{5}{4^{r}}\right)+O\left(n^{-1}\right), \quad \mathbb{V}=\frac{4^{r}-1}{3 \cdot 16^{r}} n+O(1)
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- $C \approx 1.36190, \delta \ldots$ periodic fluctuation


## Size of $r$-fold reduced plane trees

## Leaves

$\mathbb{E} \sim \frac{n}{r+1}$
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limit law: $\checkmark$


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## Disclaimer

Results are not always that nice!

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- Q: How many branches are there?



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Theorem (H.-Heuberger-Kropf-Prodinger)
Average \# of branches in a random plane tree of size $n$ is

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& -C=-\frac{\gamma+4 \alpha \log 2+\log 2+24 \zeta^{\prime}(-1)+2}{12 \log 2} \approx-0.11811,
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- $C=-\frac{\gamma+4 \alpha \log 2+\log 2+24 \zeta^{\prime}(-1)+2}{12 \log 2} \approx-0.11811$,
- $\delta$. . . periodic fluctuation:

$$
\delta(x):=\frac{1}{\log 2} \sum_{k \in \mathbb{Z} \backslash\{0\}}\left(-1+\chi_{k}\right) \Gamma\left(\chi_{k} / 2\right) \zeta\left(-1+\chi_{k}\right) e^{2 k \pi i x}, \quad \chi_{k}=\frac{2 \pi i k}{\log 2}
$$

## More about my thesis? © ${ }^{-}$

- Defense @ 30.5., 11:00 / E.1.05


## Asymptotic Analysis of Shape Parameters of Trees and Lattice Paths <br> PhD Thesis / Defense

