## Benjamin HackI

## Asymptotic Analysis of Shape Parameters of Trees and Lattice Paths

KARL
POPPER
KOLLEG

## Thesis Overview

(1) Reductions of Binary Trees and Lattice Paths induced by the Register Function

- Joint work with Clemens Heuberger, Helmut Prodinger
- Published: Theoret. Comput. Sci. 705 (2018), 31-57.

2) Fringe Analysis of Plane Trees Related to Cutting and Pruning

- Joint work with Clemens Heuberger, Sara Kropf, Helmut Prodinger
- Published: Aequationes Math. 92 (2018), 311-353.
(3) Growing and Destroying Catalan-Stanley Trees
- Joint work with Helmut Prodinger
- Published: Discrete Math. Theor. Comput. Sci. 20 (2018).

4. Ascents in Non-Negative Lattice Paths

- Joint work with Clemens Heuberger, Helmut Prodinger
- arXiv:1801.02996 [math.CO]


## Example: Deterministic Tree Reduction

- Remove all leaves!



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Parameters of Interest:

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- Size of $r$ th reduction


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Parameters of Interest:

- Size of $r$ th reduction
- Age: \# of possible reductions


## Reduction $\rightarrow$ Expansion

- modelling reduction directly: not suitable

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- instead: see inverse operation, growing trees
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- F...family of plane trees; bivariate generating function $f$
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$$
\square \triangleq t, \bigcirc \triangleq z \quad \Rightarrow \quad \Phi(t)=z t+z t^{2}+z t^{3}+\cdots
$$

## Reductions on Plane Trees

Leaves


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Old leaves


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## Reductions on Plane Trees



Old paths


## Reductions on Plane Trees



Paths $\bigcirc$


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## Parameters of Interest:

- tree size after $r$ reductions


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## Bivariate Generating Function

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\Rightarrow T(z, t)=\frac{1-(z-t)-\sqrt{1-2(z+t)+(z-t)^{2}}}{2}
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Proof. Symbolic equation

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\mathcal{T}=\square+\underbrace{}_{\mathcal{T}} \underbrace{}_{\mathcal{T}}
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- As $\Phi$ is linear, this proves the proposition.


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## Cutting leaves

## Theorem (H.-Heuberger-Kropf-Prodinger)

After $r$ reductions of a random tree of size $n$, the remaining size $X_{n, r}$ has mean and variance

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## Proof insights:

- $\mathbb{E} X_{n, r}$ and $\mathbb{V} X_{n, r}$ follow via singularity analysis
- Asymptotic normality: $n-X_{n, r}$ is a tree parameter with small toll function, limit law by Wagner (2015)



## Pruning

- Remove all paths that end in a leaf!



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## Branches in a Tree

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Proof: all branches end in exactly one leaf (at some point).

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Theorem (H.-Heuberger-Kropf-Prodinger)
Average \# of branches in a random plane tree of size $n$ is

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- $C=-\frac{\gamma+4 \alpha \log 2+\log 2+24 \zeta^{\prime}(-1)+2}{12 \log 2} \approx-0.11811$,
- $\delta$. . . periodic fluctuation:

$$
\delta(x):=\frac{1}{\log 2} \sum_{k \in \mathbb{Z} \backslash\{0\}}\left(-1+\chi_{k}\right) \Gamma\left(\chi_{k} / 2\right) \zeta\left(-1+\chi_{k}\right) e^{2 k \pi i x}, \quad \chi_{k}=\frac{2 \pi i k}{\log 2} .
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## Summary: Reductions on Plane Trees

## Leaves

$\mathbb{E} \sim \frac{n}{r+1}$
$\mathbb{V} \sim \frac{r(r+2)}{6(r+1)^{2}} n$
limit law: $\checkmark$


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## Disclaimer

Results are not always that nice!

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## Counterexample: Catalan-Stanley trees

- Motivation: Stanley's Catalan interpretation \#26
- Rightmost leaves in all branches of root have odd distance



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## Counterexample: Results

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## Age

## Size of $r$ th Reduction

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$\bigcirc \quad \mapsto \quad \mapsto$



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$\bigcirc \quad \mapsto \quad \cdots \quad \mapsto$



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- $\mathbb{E} \sim \frac{1}{4^{\prime}} n$
$-\mathbb{V} \sim \frac{\left(2^{r}+1\right)\left(2^{r}-1\right)}{16^{r}} n^{2}$



## Trimming Binary Trees

Cutting strategy:

- Remove Leaves
- Merge single children with their corresponding parent



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We label the nodes according to the following rules:

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- r-branches, Numerics: Yamamoto, Yamazaki (2009)


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- Number / Distribution of ( $r$ - $)$ branches?
- Example: | $r$ | 0 | 1 | 2 | 3 |
| ---: | :---: | :---: | :---: | :---: |$\#$-branches $14 \begin{array}{ll}5 & 2\end{array} 1$


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In a random binary tree of size $n .$. .

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$$
\mathbb{E}=\frac{n}{4^{r}}+\frac{1}{6}\left(1+\frac{5}{4^{r}}\right)+O\left(n^{-1}\right), \quad \mathbb{V}=\frac{4^{r}-1}{3 \cdot 16^{r}} n+O(1)
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$$

- expected total \# of branches is

$$
\frac{4}{3} n+\frac{1}{6} \log _{4} n+C+\delta\left(\log _{4} n\right)+O\left(n^{-1} \log n\right)
$$

- $C \approx 1.36190, \delta \ldots$ periodic fluctuation


## Non-Negative Lattice Paths

## Dyck Paths:

- Sequences of $\{-1,1\} \triangleq\{\searrow, \nearrow\}$,
- Never below axis, end on axis.



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Łukasiewicz Excursions:

- Sequences of $\mathcal{S}=\{-1\} \cup N, N \subseteq \mathbb{N}_{0}$,
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## Ascents



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- $r$-Ascent: ascent of length $r$.


## Bijection: Excursions $\longleftrightarrow$ Trees



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## Ascents and Generating Functions <br> - $V(z, t) \ldots$ BGF for Plane $\operatorname{Trees}(\mathcal{S}+1)$ <br> $>z \rightsquigarrow$ tree size, $t \rightsquigarrow \#$ of $r$-ascents

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## Proposition

- $\tau \ldots$...structural constant, unique $\tau>0$ with $S^{\prime}(\tau)=0$
- $p=\operatorname{gcd}(\mathcal{S}+1) \ldots$ period, $\zeta^{p}=1$
$V(z)$ has radius of convergence $\rho=1 / S(\tau)$, singularities at $\zeta \rho$ and

$$
V(z) \stackrel{z \rightarrow \zeta \rho}{=} \zeta \tau-\zeta \sqrt{\frac{2 S(\tau)}{S^{\prime \prime}(\tau)}}\left(1-\frac{z}{\zeta \rho}\right)^{1 / 2}+O\left(1-\frac{z}{\zeta \rho}\right)
$$

## Analysis of $r$-Ascents in Excursions

Theorem (H.-Heuberger-Prodinger)

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$$
\begin{aligned}
\mathbb{V}=\left(\frac{(c-1)^{r}}{c^{r+2}}+\right. & \frac{(2 c-2 r-3)(c-1)^{2 r}}{c^{2 r+4}} \\
& \left.-\frac{(c-1)^{2 r-2}(2 c-r-2)^{2}}{c^{2 r+3} \tau^{3} S^{\prime \prime}(\tau)}\right) n+O\left(n^{1 / 2}\right)
\end{aligned}
$$

## Excursions - Example

## Example ( $r$-Ascents in Dyck paths)

- $\mathcal{S}=\{-1,1\}, p=2, \tau=1$.


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\begin{aligned}
& \mathbb{E} D_{2 n, r}=\frac{n}{2^{r+1}}-\frac{(r+1)(r-4)}{2^{r+3}} \\
& \quad+\frac{\left(r^{2}-11 r+22\right)(r+1) r}{2^{r+6}} n^{-1}+O\left(n^{-2}\right)
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& \mathbb{V} D_{2 n, r}=\left(\frac{1}{2^{r+1}}-\frac{r^{2}-2 r+3}{2^{2 r+3}}\right) n+O(1)
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$$

## Summary <br> Bijection



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## Lagrange Inversion

- $y, \Phi, H \ldots$ formal power series with $y=x \Phi(y)$ and $\Phi(0) \neq 0$

Then:

$$
\left[x^{n}\right] H(y)=\frac{1}{n}\left[y^{n-1}\right] H^{\prime}(y) \Phi(y)^{n}
$$

## Singularity Analysis - Standard Scale

- $\alpha \in \mathbb{C} \backslash \mathbb{Z}_{\leq 0}$
- $f(z)=(1-z)^{-\alpha}$

Then:

$$
\left[z^{n}\right] f(z) \sim \frac{n^{\alpha-1}}{\Gamma(\alpha)}\left(1+\frac{e_{1}(\alpha)}{n}+\frac{e_{2}(\alpha)}{n^{2}}+\ldots\right),
$$

where

$$
\begin{aligned}
& e_{1}(\alpha)=\frac{\alpha(\alpha-1)}{2} \\
& e_{2}(\alpha)=\frac{\alpha(\alpha-1)(\alpha-2)(3 \alpha-1)}{24} \\
& e_{3}(\alpha)=\frac{\alpha^{2}(\alpha-1)^{2}(\alpha-2)(\alpha-3)}{48} .
\end{aligned}
$$

## Singularity Analysis - Logarithmic Scale

- $\alpha \in \mathbb{C} \backslash \mathbb{Z}_{\leq 0}$
- $f(z)=(1-z)^{-\alpha}\left(\frac{1}{z} \log \frac{1}{1-z}\right)^{\beta}$

Then:

$$
\left[z^{n}\right] f(z) \sim \frac{n^{\alpha-1}}{\Gamma(\alpha)}\left(1+\frac{C_{1}}{\log n}+\frac{C_{2}}{\log ^{2} n}+\ldots\right)
$$

where

$$
C_{k}=\left.\binom{\beta}{k} \Gamma(\alpha) \frac{d^{k}}{d s^{k}} \frac{1}{\Gamma(s)}\right|_{s=\alpha}
$$

## Mellin Transform - Properties

$$
\mathcal{M}\left(e^{-x}\right)(s)=\Gamma(s), \quad \mathcal{M}\left(e^{-x^{2}}\right)(s)=\frac{1}{2} \Gamma\left(\frac{s}{2}\right), \quad \mathcal{M}(\llbracket 0 \leq s \leq 1 \rrbracket)(s)=\frac{1}{s}
$$

$$
\mathcal{M}(\log (1+x))(s)=\frac{\pi}{s \sin (\pi s)}, \quad \mathcal{M}\left(\frac{1}{e^{x}-1}\right)=\zeta(s) \Gamma(s)
$$

$$
\begin{aligned}
& f^{*}(s)=\int_{0}^{\infty} f(x) x^{s-1} d x \quad f(x)=\frac{1}{2 \pi i} \int_{c-i \infty}^{c+i \infty} f^{*}(s) x^{-s} d s \\
& \text { lin. comb. lin. comb. } \\
& x^{k} f(x) \quad f^{*}(s+k) \quad\langle\alpha-k, \beta-k\rangle \\
& f\left(x^{\rho}\right) \quad \frac{1}{\rho} f^{*}\left(\frac{s}{\rho}\right) \quad\langle\rho \alpha, \rho \beta\rangle \\
& f(\mu x) \quad \mu^{-s} f^{*}(s) \quad\langle\alpha, \beta\rangle(\mu>0) \\
& f(x) \log x \quad \frac{d}{d s} f^{*}(s) \quad\langle\alpha, \beta\rangle \\
& x \frac{d}{d x} f(x) \quad-s f^{*}(s)
\end{aligned}
$$

## Mellin Transform - Correspondence

| $f(x)$ | $f^{*}(s)$ |
| :--- | :--- |
| $f(x)=O\left(x^{\alpha}\right), x \rightarrow 0$ | $-\alpha$ left border of FS |
| $f(x)=O\left(x^{\beta}\right), x \rightarrow \infty$ | $-\beta$ right border of FS |
| Expansion up to $O\left(x^{\gamma}\right), x \rightarrow 0$ | Meromorph. Cont. up to Re $s>-\gamma$ |
| Expansion up to $O\left(x^{\delta}\right), x \rightarrow \infty$ | Meromorph. Cont. up to Re $s<-\delta$ |
| Growth $x^{k} \log ^{\ell} x, x \rightarrow 0$ | Pole with expansion $\frac{(-1)^{\ell} \ell!}{(s+k)^{\ell+1}}$ |
| Growth $x^{k} \log ^{\ell} x, x \rightarrow \infty$ | Pole with expansion $-\frac{(-1)^{\ell} \ell!}{(s+k)^{\ell+1}}$ |

## Hurwitz Zeta Function - Estimate

$$
\zeta(s, \alpha):=\sum_{n>-\alpha} \frac{1}{(n+\alpha)^{s}}, \quad \operatorname{Re} s>1
$$

- $s=\sigma+i t, \sigma_{0} \leq \sigma \leq \sigma_{1}$
- $|t| \rightarrow \infty$

$$
|\zeta(s, \alpha)|=O\left(|t|^{\tau(\sigma)} \log |t|\right), \quad \tau(\sigma)= \begin{cases}\frac{1}{2}-\sigma, & \sigma \leq 0 \\ 1 / 2, & 0 \leq \sigma \leq \frac{1}{2} \\ 1-\sigma, & \frac{1}{2} \leq \sigma \leq 1 \\ 0, & \sigma \geq 1\end{cases}
$$

The logarithmic factor is only necessary in a neighborhood of $\sigma=0, \sigma=1$.

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## Gamma Function - Estimate

$$
\Gamma(z):=\int_{0}^{\infty} e^{-t} t^{z-1} d t
$$

- $z \rightarrow \infty$ with $|\arg z| \leq \pi-\delta, \delta>0$

$$
\Gamma(z) \sim e^{-z} z^{z}\left(\frac{2 \pi}{z}\right)^{1 / 2}\left(1+\frac{1}{12 z}+\frac{1}{288 z^{2}}+\ldots\right)
$$

